A HYBRID SWARM INTELLIGENCE BASED PARTICLE BEE ALGORITHM FOR BENCHMARK FUNCTIONS AND CONSTRUCTION SITE LAYOUT OPTIMIZATION

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ABSTRACT: The construction site layout (CSL) design presents a particularly interesting area of study because of its relatively high level of attention to usability qualities, in addition to common engineering objectives such as cost and performance. However, it is difficult combinatorial optimization problem for engineers. Swarm intelligence (SI) was very popular and widely used in many complex optimization problems which was collective behavior of social systems such as honey bees (bee algorithm, BA) and birds (particle swarm optimization, PSO). In order to integrate BA global search ability with the local search advantages of PSO, this study proposes a new optimization hybrid swarm algorithm – the particle bee algorithm (PBA) which imitates the intelligent swarming behavior of honeybees and birds. This study compares the performance of PBA with that of genetic algorithm (GA), differential evolution (DE), bee algorithm (BA) and particle swarm optimization engineering of CSL problem. The results show that the performance of PBA with that of BA and PSO for practical construction engineering of CSL problem. The results show that the performance of PBA is comparable to those of the mentioned algorithms in the benchmark functions and can be efficiently employed to solve a hypothetical CSL problem with high dimensionality.

Keywords: Construction Site Layout, Swarm Intelligence, Bee Algorithm, Particle Swarm Optimization, Particle Bee Algorithm

1. INTRODUCTION

Construction site layout (CSL) problems are particularly interesting because in addition to common engineering objectives such as cost and performance, facility design is especially concerned with aesthetics and usability qualities of a layout [1]. The CSL problem identifies a feasible location for a set of interrelated objects that meet all design requirements and maximizes design quality in terms of design preferences while minimizing total cost associated with interactions between these facilities. Pairwise costs usually reflect transportation costs and/or inter-facility adjacency preferences [1, 2]. CSL problems arise in the design of hospitals, service centers and other facilities [3]. However, all such problems are known as "NP-hard" and because of the combinatorial complexity, it cannot be solved exhaustively for reasonably sized layout problems [3].

In the past, Elbeitagi and Hegazy [4] used a hybrid neural network to identify optimal site layout. Yeh [3] applied annealed neural networks to solve construction site-level CSL problems. Gero and Kazakov [5] incorporated the concept of genetic engineering into the GA system for solving building space layout problems. Li and Love [6] and Osman et al. [7] used GA to solve site layout problems in unequally sized facilities. The objective functions of the above-mentioned algorithms were to optimize the interaction between facilities, such as total inter-facility transportation costs and frequency of inter-facility trips. Those previous research focused on solving different optimization problems by applying those algorithms under different constraints which quality of solutions were limited by the capability of the algorithms.

Swarm intelligence (SI) has been of increasing interest to research scientists in recent years. SI was defined by Bonabeau et al. as any attempt to design algorithms or distributed problem-solving devices based on the collective behavior of social insect colonies or other animals [8]. Bonabeau et al. focused primarily on the social behavior of ants [9], fish [10], birds [11] and bees [12] etc. However, the term "swarm" can be applied more generally to refer to any restrained collection of interacting agents or individuals. Although bees swarming around a hive is the classical example of "swarm", swarms can easily be extended to other systems with similar architectures.

A few models have been developed to model the intelligent behaviors of honeybee swarms and applied to solve combinatorial type problems. Karaboga et al. [13] presented an artificial bee colony (ABC) algorithm and expanded its experimental results [14]. It has been pointed out that the ABC algorithm outperforms GA for functions exhibiting multi-modality or uni-modality. Pham et al. [12] presented an original bee algorithm (BA) and applied to two standard functional optimization problems with two and six dimensions. Results demonstrated BA able to find solutions very close to the optimum, showing that BA generally outperformed GA. However, while BA [12] offers the potential to conduct global searches and uses a simpler mechanism in comparison with GA, it dependence on random search makes it relatively weak in local search activities and does not records past searching experiences during the optimization search process. For instance, a flock of birds may be thought of as a swarm whose individual agents are birds. Particle swarm optimization (PSO), which has become quite popular for many researchers recently, models the social behavior of birds [12]. PSO potentially used in local searching, and records past searching experiences during optimization search process. However, it converges early in highly discrete problems [15].

Hence, in order to improve BA and PSO, this study proposed an improved optimization hybrid swarm algorithm called the particle bee algorithm (PBA) that imitates a particular intelligent behavior of bird and honey bee swarms and integrates their advantages. In addition, this study also proposed a neighborhood-windows (NW) technique for improving PBA search efficiency and proposed a self-parameter-updating (SPU) technique for preventing trapping into a local optimum in high dimensional problems. This study compares the performance of PBA algorithm with that of BA [12] and PSO for a hypothetical construction engineering of CSL problem.

2. HYBRID SWARM PARTICLE BEE ALGORITHM

For improved BA local search ability, PSO global search ability and to seek records past experience during optimization search process, this study reconfigures the neighborhood dance search [12] as a PSO search [11]. Based on cooperation between bees (BA) and birds (PSO), the proposed algorithm improves BA neighborhood search using PSO search. Therefore, PBA employs no recruit bee searching around "elite" or "best" positions (as BA does). Instead, a PSO search is used for all elite and best bees. In other words, after PSO search, the number of "elite", "best" and "random" bees equals the number of scout bees. In PBA, the particle bee colony contains four groups, namely (1) number of scout bees (n), (2) number of elite sites selected out of n visited sites (e), (3) number of best sites out of n visited sites (b), and (4) number of bees recruited for the other visited sites (r). The first half of the bee colony consists of elite bees, and the second half includes the best and random bees. The particle bee colony contains two parameters, i.e., number of iteration for elite bees by PSO (Peitr) and number of iteration for best bees by PSO (Pbitr). PBA flowchart is shown in Fig. 1.



Step (1) Initialize scout bees: PBA starts with *n* scout bees being randomly placed with respective positions and velocities in the search space.

Step (2) Evaluate fitness: Start the loop and evaluate scout bee fitness.

Step (3) Select elite sites (*e*) **from scout bees:** Elite sites are selected for each elite bee, whose total number is equal to half the number of scout bees.

Step (4) Elite bees initiate the PSO procedure by *Peitr* iteration for neighborhood-windows (NW): In this step, new particle bees from elite and best bees are produced using Eq. (1). Elite and best bee velocity updates are performed as indicated in Eq. (2). This study further proposed a neighborhood-windows (NW) technique to improve PSO searching efficiency as show in Eq. (3). Thus, after $x_{id}(t+1)$ is substituted into Eq. (1) and Eq. (2), the NW ensures PSO searching within the designated x_{idmin} and x_{idmax} . In other word, if the sum of $x_{id}(t+1)$ exceeds x_{idmin} or x_{idmax} then $x_{id}(t+1)$ is limited to x_{idmin} or x_{idmax} .

 $x_{id}(t+1) = x_{id}(t) + v_{id}(t+1)$(1) where x_i is $i_{th} x$ and i = 1 to n; v_i is $i_{th} v$; d is dimension in x_i or v and d = 1to D; t is iteration; $x_{id}(t)$ is d_{th} dimension in $i_{th} x$ and in t iteration; $v_{id}(t+1)$ is d_{th} dimension in $i_{th} v$ and in t+1 iteration; $x_{id}(t+1)$ is d_{th} dimension in $i_{th} x$ and in t+1 iteration; n is number of particles.

 $x_{id \min} \le x_{id} (t+1) \le x_{id \max}$(3) where x_i is $i_{th} x$ and i = 1 to n; d is dimension in x_i and d = 1 to D; t is iteration; $x_{id}(t+1)$ is d_{th} dimension in $i_{th} x$ and in t+1 iteration; n is number of particles.

distributed real random number within the range 0 to 1.

Step (5) Select best sites (b) from scout bees: Best sites are selected for each best bee, the total number of which equals one-quarter of the number of scout bees.

Step (6) Best bees start the PSO procedure using the NW *Pbitr* iteration: In this step, new particle bees from elite and best bees are produced using **Eq.** (1). Elite and best bee velocity updates are acquired using **Eq.** (2). The

NW technique improves PSO search efficiency, as show in **Eq. (3)**.

Step (7) Recruit random bees (r) for other visited sites: The random bees in the population are assigned randomly around the search space scouting for new potential solutions. The total number of random bees is one-quarter of the number of scout bees.

Step (8) Self-parameter-updating (SPU) for elite, best and random bees: Furthermore, in order to prevent being trapped into a local optimum in high dimensional problems, this study proposed a solution, i.e., the self-parameterupdating (SPU) technique, the idea for which came from Karaboga [13]. Eq. (4) shows the SPU equation.

$x_{id}(new) = x_{id}(cur) + 2 \times (\text{Rand-0.5}) \times (x_{id}(cur) - x_{jk}) \dots \dots$
$j = int (Rand \times n) + 1(5)$
$k = \operatorname{int} (\operatorname{Rand} \times d) + 1(6)$
where x_i is $i_{th} x$ and $i = 1$ to n ; d is dimension in x_i and $d = 1$ to D ; $x_{id}(cur)$
is d_{th} dimension in $i_{th} x$ and in current solution; $x_{id}(new)$ is d_{th} dimension in
ith x and in new solution; Rand is a uniformly distributed real random
number within the range 0 to 1; j is the index of the solution chosen
randomly from the colony as shows in Eq. (5) , k is the index of the
dimension chosen randomly from the dimension as shows in Eq. (6) ; n is
number of scout bees.

In step (8), after elite, best and random bees have been distributed based on finesse, finesses are checked to determine whether they are to be abandoned or memorized using Eq. (4). Therefore, if finesses of elite, best or random bees are both improved using Eq. (4) and improved over previous finesses, the new finesses are memorized. In step (3) through step (8), this differential recruitment is a key operation of the PBA. The Rastrigin function is provided as a sample for observing the proposed optimization procedure behavior in this study. Rastrigin function formula and figure are shown, respectively, in Eq. (7). In this study, while the NW carries out the PSO search of the local optimization at the start of each iteration search space, BA random bees and the SPU technique control the global optimization at the end of each iteration search space. Therefore, in the PBA procedure, the proposed optimization search techniques may search across each iteration search space based on that particular search space's potential. The procedure of this study combines

NW, BA's random bees and the SPU technique, as show in **Fig. 2** and **Fig. 3**.

 $f(\vec{x}) = \sum_{d=1}^{2} (x_{id}^2 - 10\cos(2\pi x_{id}) + 10)$ where -5.12 x_{id} 5.12; f(0) = 0(7)



Fig.2 Optimization procedure behavior in two dimensions



Fig.3 Optimization procedure behavior in each iteration

Step (9) Convergence?: In this step, only the bee with the highest fitness will be selected to form the next bee population. These steps are repeated until the stop criterion is met and bees are selected to be abandoned or memorized.

3. BENCHMARK FUNCTIONS EXPERIMENT

3.1 Modeling of Benchmark functions

Some classical benchmark functions used by Karaboga [13] are presented, in order to evaluate PBA performance. This study compares PBA results against the GA, DE and PSO results of Karaboga [13] and equalizes the total number of evaluations to 500,000 for all the functions, as shown in Ref. [13]. In PBA, maximum number of cycles was designated for *Bitr, Peitr, Pbitr* as, respectively, 5,000, 15, 9 for all the functions. The number of elite bees totaled 50 percent of the colony and the number of best bees totaled

25 percent of the colony. In PBA, the number of random bees is taken *n-e-b* from the number of elite bees and best bees. The PBA used in the simulation studies and values assigned for the parameter settings of GA, DE, PSO in Ref. [13] and BA in the Ref. [12] are given in Table 1.

 Table 1
 Parameter values used in the experiments

GA	GA[13] DE[13]		PSO		BA		PBA		
n	50	п	50	п	50	п	50	п	50
m	0.01	С	0.9	W	0.9~0.7	е	n/2	е	n/2
С	0.8	f	0.5	v	$X_{min}/10 \sim X_{max}/10$	b	n/4	b	n/4
g	0.9					r	n/4	r	n/4
						n_I	2	w	0.9~0.7
						n_2	1	v	$X_{min}/10 \sim X_{max}/10$
								Peitr	15
								Pbitr	9

n=population size (colony size); *m*=mutation rate; *c*=crossover rate; *g*=generation gap; *f*=scaling factor; *w*=inertia weight; *v*=limit of velocity; *e*=elite bee number; *b*=best bee number; *r*=random bee number; *n*1= elite bee neighborhood number; *n*2=best bee neighborhood number; *Peitr*=PSO iteration of elite bees; *Pbitr*=PSO iteration of best bees.

3.2 Results and discussion for benchmark functions

Each experiment ran for 30 runs, and average function values for the best solutions were found and recorded. Mean and standard deviations of function values obtained by DE [13], EA [13], PSO [13], BA [12] and PBA under the same conditions are given in Table 1. In order to compare the performance of those algorithm, this study focus on the performance of "algorithm found global optimum" and "solution better than others" as shown in Table 2. From Table 2, the total score of PBA was 37 which better than BA, DE, PSO and GA respectively 33, 32, 30 and 13. In Table 3, the algorithm performance cross-matching on benchmark functions is presented. As shown in Table 3, PBA has better performance then others. From results presented in Table 2 and Table 3, PBA has better performance amongst all algorithms considered in the present investigation.

 Table 2
 Algorithm performance comparison

	GA	DE	PSO	BA	PBA
Count of algorithm found global optimum	9	18	17	18	20
Count of solution better than others	4	14	13	15	17
Total score	13	32	30	33	37

,	Table	3 Algor	rithm pe	rforman	ce cross-	matching
		GA	DE	PSO	BA	PBA
	GA	(0, 0, 0)	(1, 16, 9)	(3, 15, 8)	(4, 16, 6)	(0, 16, 10)
	DE	(16, 1, 9)	(0, 0, 0)	(6, 3, 17)	(6, 5, 15)	(4, 5, 17)
	PSO	(15, 3, 8)	(3, 6, 17)	(0, 0, 0)	(6, 7, 13)	(3, 6, 17)
	BA	(16, 4, 6)	(5, 6, 15)	(7, 6, 13)	(0, 0, 0)	(4, 6, 16)
	DD A	(1(0, 1, 0, 1, 0))	(5 4 17)	((2 17)	((1 1()	(0, 0, 0)

Notice: The contents in bracket are Win, Lose and Fair. Win/Lose/Fair means the algorithm has the better/worse/fair solutions with the compared algorithm.

4. CONSTRUCTION SITE LAYOUT PROBLEM

4.1 Modeling of a construction site layout problem

4.1.1 A hypothetical construction site layout

A medium-sized project is adopted as a hypothetical construction site **[16]** to determine optimal site layout through PSO, BA and PBA.

4.1.2 Number and type of facilities

This study considers some common site facilities, such as a site office, a labor hut, a materials storage area, a main gate and a refuse storage area [16]. The numbered site facilities are listed in Table 4.

Table 4	Facilities used on the cas	se study
Facility no.	Facility name	Note
A	Site office	-
в	Debris storage area	-
C	Reinforcement bending/storage yard	

C	Reinforcement bending/storage yard	-
D	Carpentry workshop and store	-
E	Labor hut	-
F	Materials storage area	-
G	Main gate (fixed)	Fixed
Н	Materials hoist (fixed)	Fixed
Ι	Refuse chute (fixed)	Fixed

4.1.3 Travel distance between site locations

The travel distance between locations is measured using the rectangular distance representing the actual operations and resource movements on site. **Table 5** shows the travel distances between the possible facility locations **[16]**.

Dicto		Location												
Dista	nce	1	2	3	4	5	6	7	8	9	10	11	12	13
	1	0	1	2	6	7	9	12	14	16	17	13	4	9
	2	1	0	1	5	6	8	11	13	15	16	12	3	8
	3	2	1	0	4	5	7	10	12	14	15	11	2	7
	4	6	5	4	0	1	3	7	9	11	12	9	2	5
	5	7	6	5	1	0	2	6	8	10	11	8	3	4
Lo	6	9	8	7	3	2	0	3	5	7	8	8	5	4
cati	7	12	11	10	7	6	3	0	2	4	5	7	6	3
ion	8	14	13	12	9	8	5	2	0	2	3	5	8	3
	9	16	15	14	11	10	7	4	2	0	1	3	11	6
	10	17	16	15	12	11	8	5	3	1	0	2	12	7
	11	13	12	11	9	8	8	7	5	3	2	0	9	5
	12	4	3	2	2	3	5	6	8	11	12	9	0	4
	13	9	8	7	5	4	4	3	3	6	7	5	4	0

4.1.4 Trip frequency between facilities

Trip frequency between facilities influences site layout planning and the proximity of predetermined site facilities. Therefore, the frequency of trips made between facilities in a single day are assumed **[16]** as shown in **Table 6**.

4.1.5 Objective function

This study is based on Ref. [16] that gives the total objective function as follows Eq. (8):

Minimize
$$\sum_{i=1}^{n} \sum_{j=1}^{n} \sum_{k=1}^{n} x_{ij} \times 1/f_{ik} \times d_{ik}$$
(8)

Subject to

If no reasonable solutions exist that the value of results is 150

If i=j then $x_{ij}=0$; If i=k then $f_{ik}=d_{ik}=0$

where *n* is the number of facilities; x_{ij} is the permutation matrix variable such that when facility *i* is assigned to location *j*; f_{ik} is the proximity relationship between facilities *i* and *k*; and d_{ik} is the distance between locations *i* and *k*.

Fable	6	Frea	uencies	of tri	ns bet	ween	faci	lities
	v			01 011	P0 00		10001	

Frequency			Facility										
riequ	che y	Α	В	С	D	E	F	G	Н	Ι			
	Α	0	3.11	4.79	4.94	5.15	5.41	6.34	3.48	2.55			
	В	3.11	0	3.69	3.71	3.7	3.36	4.42	3.07	5.85			
	С	4.79	3.69	0	4.27	4	4.4	5.65	6.26	2.03			
Ŧ	D	4.94	3.71	4.27	0	4.51	4.58	5.14	6.2	2.24			
cil	Е	5.15	3.7	4	4.51	0	4.99	4.39	4.13	2.48			
ity	F	5.41	3.36	4.4	4.58	4.99	0	5.24	6.2	2.65			
	G	6.34	4.42	5.65	5.14	4.39	5.24	0	4.62	3.75			
	Н	3.48	3.07	6.26	6.2	4.13	6.2	4.62	0	2.37			
	Ι	2.55	5.85	2.03	2.24	2.48	2.65	3.75	2.37	0			

4.2 Results and discussion for CSL problem

This study was adapted from 30 experimental runs with the values found in **Table 1** through 500 iterations by BA, PSO and PBA. **Table 7** present the evolution of the CSL problem result. As seen in **Table 7**, the best mean and lowest standard error for PBA are respectively 108.37 and 0.0389, which is better than BA (112.83 and 7.3584) and PSO (126.19 and 18.626). Besides, although BA and PSO both obtain the same best value as PBA (108.36) neither can avoid unreasonable solution results with the worst value (i.e., 150). Therefore, PBA provides a better evolution of mean fitness, standard, and best and worst result that than BA and PSO.

	Table 7	The result o		
	Mean	Std	Best	Worst
BA	112.83	7.3584	108.36	150.00
PSO	126.19	18.626	108.36	150.00
PBA	108.37	0.0389	108.36	108.47

The best layout alternatives of PBA and Lam [16] are shown in Figs. 4 and Fig. 5. The best layout design of PBA and Lam [16] is 108.36 and 114.3, respectively, as calculated by Eq. (8). Figs. 4 and Fig. 5 both place the site office near the main gate so that site staff can enter the site office via the shortest route. In a practical construction job site, the labor hut should be adjacent to the site office so that the residential area for the site staff and workers can be concentrated in a particular zone, and so the construction plan is easy to navigate. The PBA result places the site office is near labor hut (**Fig. 4**). Furthermore, it is potentially dangerous for the site manager/staff to have to travel from the site office to the labor hut through the debris storage area, the materials storage area, and the carpentry workshop and store (see **Fig. 5**). Besides, the short distance between the materials hoist and materials store (for which PBA is better than Ref. [16]) means that site workers can efficiently transport materials to the superstructure. Thus, PBA results for this hypothetical CSL problem may better than Ref. [16].





Fig. 5 Lam [16] best layout design

5. CONCLUSION

In the previous section, the performance of the particle bee algorithm (PBA) was compared with genetic algorithm (GA), differential evolution (DE), particle swarm optimization (PSO), and bee algorithm (BA) in terms of both multi-dimensional and multimodal numeric problems. Results show that PBA performs better than the mentioned algorithms on each benchmark numerical function. In the hypothetical of construction site layout (CSL) problem, the evolution of mean and best fitness, PBA are 108.37 and 0.03888 better than BA are 112.83 and 7.3584 and PSO are 126.19 and 18.626. Besides, in the comparison between PBA and reference, the results show that the PBA are more reasonable than reference. The result shows that PBA performs better than the mentioned algorithms for this hypothetical CSL problem.

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